

# Solution analysis and continuation algorithms for multilinear pagerank

Federico Poloni  
Joint work with B. Meini

University of Pisa

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# Pagerank extensions

## Pagerank [Page '98]

**Input:** transition probabilities  $P_{ij} = P[i \rightarrow j]$ , ‘personalization vector’  $v$ .

**Output:** ‘importance’ score  $x_i$  of each node

$$x = \alpha Px + (1 - \alpha)v, \quad \mathbf{1}^\top x = 1.$$

## Multilinear Pagerank [Gleich-Lim-Yu '15]

**Input:** **two-step** transition probabilities  $P_{ijk} = P[i \rightarrow j \rightarrow k]$ ,  
‘personalization vector’  $v$ .

**Output:** ‘importance’ score  $x_i$  of each node

$$x = \alpha \sum_{i,j=1}^n P_{ij} \cancel{x_i x_j} + (1 - \alpha)v, \quad \mathbf{1}^\top x = 1.$$

**Not** just a second-order Markov chain — that would rank edges,  $X_{ij}$ .

## Ways to look at it

Viewpoint 1: tensor eigenvalues.

$$x = \alpha \sum_{i,j=1}^n P_{ij} x_i x_j + (1 - \alpha) v \mathbf{x}_i \mathbf{x}_j, \quad \mathbf{1}^\top x = 1.$$

Since  $\mathbf{1}^\top x = 1$ , we can homogenize. Z-eigenvector problem.

[e.g., Qi, Chen, Chen book '18]

Viewpoint 2: branching processes and quadratic vector equations.

$$x = \sum_{i,j=1}^n \hat{P}_{ij} x_i x_j + \hat{v}, \quad x \geq 0.$$

[Kolmogorov 1940s, Etessami-Yannakakis '05, Bean-Kontoleon-Taylor '08, P. '13]

We focus on Viewpoint 2 here.

## Quadratic vector equations: the theory

$$\hat{P} := \alpha P, \hat{v} := (1 - \alpha)v.$$

$$x = \sum_{i,j=1}^n \hat{P}_{ij} x_i x_j + \hat{v}, \quad x \geq 0.$$

- Among all nonnegative solutions, there is a **minimal** one  $x_*$  (i.e.,  $0 \leq x_* \leq x$  for all other solutions  $x$ ).
- If we start from  $x^{(0)} = 0$ , most iterations (fixed-point, Newton, ...) converge **monotonically** to  $x_*$ .
- $x_*$  is typically **easy** to compute, but algorithms are slower when there are other non-minimal solutions very close to it. Double-trouble.

[Hautphenne-Latouche-Remiche '11, Etessami-Stewart-Yannakakis '12, P '13]

## Adapting the theory

In MLPR,  $\mathbf{1}^\top P = \mathbf{1}^\top$ ,  $\mathbf{1}^\top v = 1$ . The iteration

$$x^{(k+1)} = \sum_{i,j=1}^n \hat{P}_{ij} x_i^{(k)} x_j^{(k)} + \hat{v}, \quad x^{(0)} = \mathbf{0}$$

has predictable behaviour on the ‘total mass’: if  $\mathbf{1}^\top x^{(k)} = w$ , then  $\mathbf{1}^\top x^{(k+1)} = \alpha w^2 + (1 - \alpha)$ .

### Theorem

- If  $\alpha \leq \frac{1}{2}$ ,  $x^{(k)} \rightarrow x_*$ , the **unique minimal** solution with  $\mathbf{1}^\top x_* = 1$ .
- If  $\alpha > \frac{1}{2}$ ,  $x^{(k)} \rightarrow x_*$ , the unique minimal solution with  $\mathbf{1}^\top x_* = \frac{1-\alpha}{\alpha}$ .

There may be **several solutions**  $x \geq x_*$  with  $\mathbf{1}^\top x = 1$ .

Uniqueness (or not) of stochastic solutions also in [Gleich-Lim-Yu '15].

# The Perron vector trick

## Idea

Is it possible to deflate the minimal solution  $x_*$ ?

Yes! (adapting an idea from [Meini-P '11] for a case with known  $x_* = \mathbf{1}$ ).

Let  $x = x_* + y$ . Then the equation becomes

$$y = H_y y, \quad (H_y) = \sum_j \hat{P}_{j:} + \hat{P}_{j::} + \hat{P}_{j::} y_j$$

$y$  = Perron eigenvector of the nonnegative matrix  $H_y$ .

- $y^{(k)}$  = Perron (maximal) eigenvector of  $H_{y^{(k-1)}}$ .
- Normalize  $y^{(k)}$  (s.t.  $x + y^{(k)}$  is stochastic).
- Iterate.

## Large and small $\alpha$

- **The good** When  $\alpha \leq \frac{1}{2}$ , there is a unique stochastic solution, and many fixed point iterations (e.g. Newton's method, Gauss-Seidel-like variants...) converge to it monotonically.
- **The bad** When  $\alpha > \frac{1}{2}$ , there may be multiple stochastic solutions, and convergence may be problematic — even if we enforce  $\mathbf{1}^\top x^{(k)} = 1$ .

## Large and small $\alpha$

- **The good** When  $\alpha \leq \frac{1}{2}$ , there is a unique stochastic solution, and many fixed point iterations (e.g. Newton's method, Gauss-Seidel-like variants...) converge to it monotonically.
- **The bad** When  $\alpha > \frac{1}{2}$ , there may be multiple stochastic solutions, and convergence may be problematic — even if we enforce  $\mathbf{1}^\top x^{(k)} = 1$ .
- **The ugly Perron-based** iterations can be used for the bad case  $\alpha > \frac{1}{2}$ :
  - ▶ Compute minimal (sub-stochastic) solution  $x^*$ ;
  - ▶ Change variable  $y = x - x^*$ ;
  - ▶ Interpret resulting equation as  $y = H_y y$ ;
  - ▶ Fixed-point iteration:  $y^{(k)} = \text{Perron eigenvector of } H_{y^{(k-1)}} \text{ (or variants)}$

(Another algorithm involving Perron vectors is in [Benson-Gleich-Lim '17].)

# Improvements

**Improvement 1:** use **Newton's method** on  $y - \text{perron}(H_y) = 0$ , where  $\text{perron}(\cdot)$  is the map that computes the Perron vector of a matrix.

An expression for the Jacobian can be found using eigenvector derivatives.

**Theorem** [Meini-P '17]

The Jacobian of the map  $w = \text{perron}(H_y)$  is

$$J = \alpha(w\mathbf{1}^\top - (I - H_y + w\mathbf{1}^\top)^{-1} \sum_j P_{:j,:} w_j).$$

(and [Bini-Meini-P '11] for the branching process problem.)

# Improvements

**Improvement 2** Use perturbation / continuation techniques: first solve the problem for an ‘easy’  $\alpha$ , then increase its value.

**Theorem** [Meini-P '17]

Let  $x_\alpha$  be the solution vector for a certain value of  $\alpha \in (0, 1)$ . Then,

$$x_{\alpha+h} = x_\alpha + \left( I - \alpha \left( \sum_j P_{j:j} + P_{j::} \right) \right)^{-1} \left( \sum_{i,j} P_{ij} (x_\alpha)_i (x_\alpha)_j - v \right) h + O(h^2).$$

Step-size heuristic: estimate the neglected second-order term  $\frac{dx_\alpha}{d\alpha^2}$ , and use it to choose the next step size.

## Variants

- Second-order Taylor
- Linear extrapolation from previous two values
- “semi-implicit method”: replace  $\alpha$  in the formula with  $\alpha + h$ .

# Experiments

For now, just **small scale examples** (but with an eye to scalability).

[Gleich-Lim-Yu '15] contains 29 small-size benchmark problems ( $n \in \{3, 4, 6\}$ ), some of them with difficult convergence.

For  $\alpha = 0.99$ , the best existing methods (newton and innout) can solve 23 and 26 of them, respectively.

**Goal:** a numerical method that scores 29/29 (in a reasonable number of iterations).

**Example (R6\_3, [Gleich-Lim-Yu '15])**

$$\frac{1}{4} \left[ \begin{array}{cccc|cccc|cccc|cccc|cccc|cccc} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 4 & 0 & 4 & 0 & 0 & 2 & 4 & 0 & 0 & 0 & 0 & 4 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 0 & 0 & 2 & 0 & 0 & 0 & 0 \\ 0 & 4 & 0 & 0 & 0 & 2 & 0 & 4 & 2 \\ 0 & 4 & 0 & 4 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 0 \\ 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 4 & 0 & 0 & 0 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 0 & 1 & 0 & 4 & 4 & 4 & 4 & 4 & 4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array} \right]$$

# Numerical results

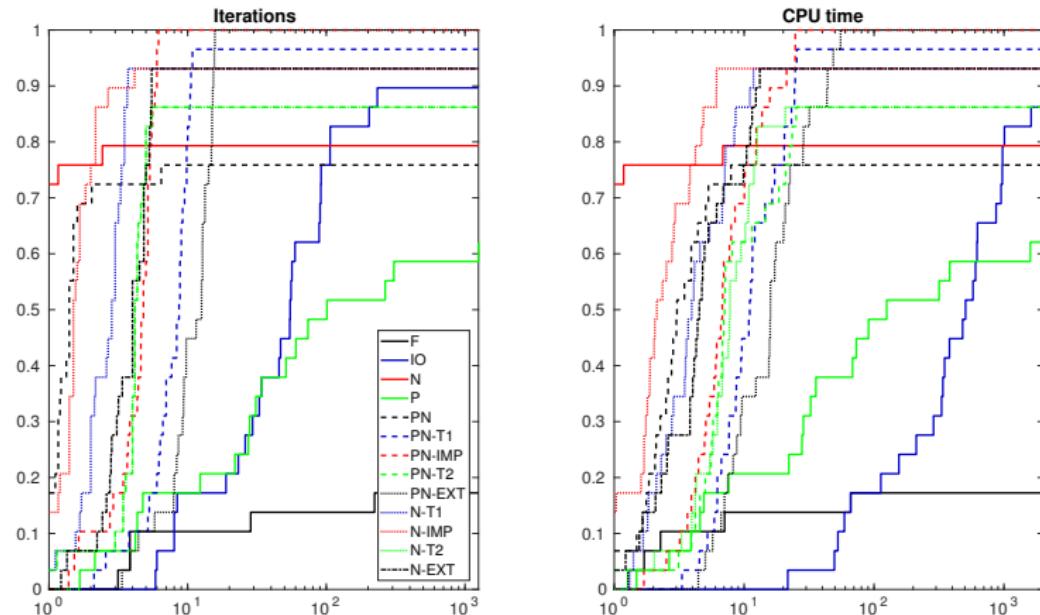


Figure: Performance profiles for the 29 benchmark tensors and  $\alpha = 0.99$

Left = faster. Top = more reliable.

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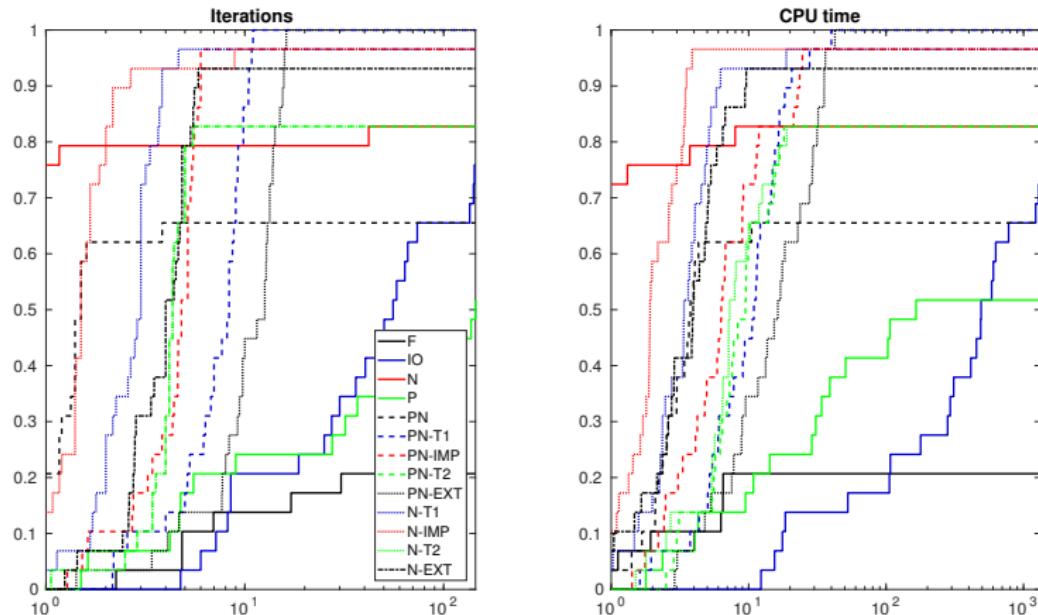


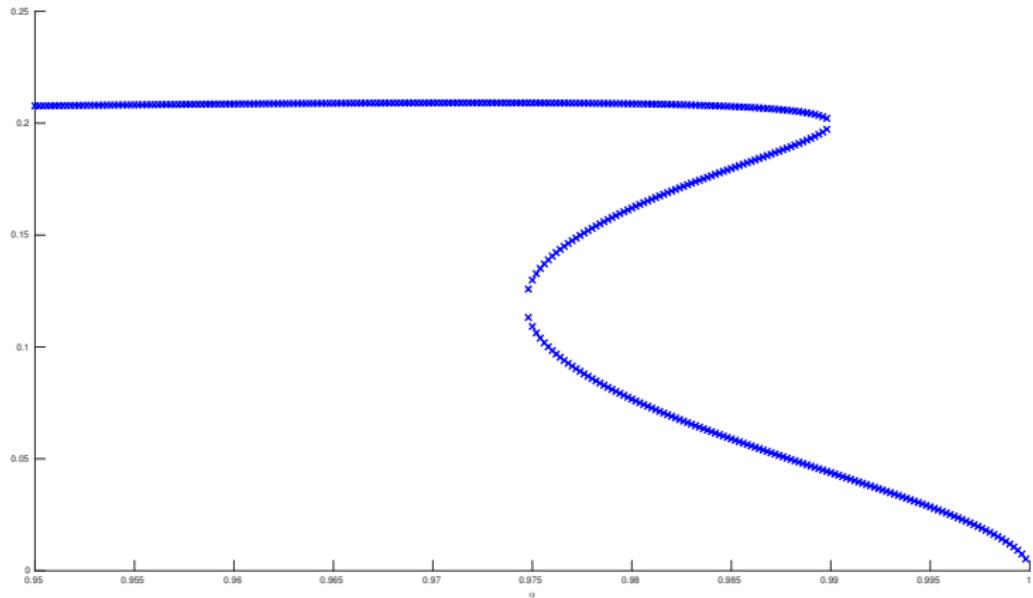
Figure: Performance profiles for the 29 benchmark tensors and  $\alpha = 0.999$

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# A difficult experiment

| F     | IO    | N   | O     | ON    | ON-T1 | ON-IMP | ON-T2 | ON-EXT | N-T1 | N-IMP | N-T2 | N-EX |
|-------|-------|-----|-------|-------|-------|--------|-------|--------|------|-------|------|------|
| NaN   | 42    | 5   | 21    | 6     | 35    | 22     | 22    | 46     | 10   | 6     | 22   | 13   |
| 23    | 48    | 6   | 26    | 7     | 50    | 36     | 24    | 80     | 17   | 13    | 24   | 29   |
| 23    | 48    | 6   | 10    | 7     | 50    | 36     | 24    | 80     | 17   | 13    | 24   | 29   |
| 114   | 24    | 4   | 19    | 6     | 28    | 22     | 22    | 39     | 8    | 6     | 22   | 11   |
| 22    | 41    | 17  | 15    | 7     | 74    | 26     | 24    | 110    | 23   | 29    | 24   | 37   |
| NaN   | 639   | 7   | NaN   | NaN   | 73    | 39     | 25    | 109    | 26   | 14    | 25   | 38   |
| NaN   | 234   | 5   | 369   | 8     | 44    | 24     | 24    | 63     | 13   | 7     | 24   | 20   |
| NaN   | 1,221 | 6   | 1,845 | 9     | 46    | 26     | 26    | 69     | 16   | 9     | 26   | 24   |
| NaN   | 532   | 5   | NaN   | 7     | 44    | 24     | 25    | 63     | 15   | 7     | 25   | 20   |
| NaN   | 339   | 6   | NaN   | 7     | 59    | 30     | 25    | 85     | 20   | 9     | 25   | 29   |
| NaN   | 228   | 5   | NaN   | 7     | 45    | 23     | 25    | 64     | 15   | 7     | 25   | 20   |
| NaN   | 357   | 6   | 362   | NaN   | 62    | 31     | 26    | 90     | 21   | 9     | 26   | 30   |
| NaN   | 156   | 6   | 7,509 | 6     | 52    | 31     | 25    | 77     | 18   | 10    | 25   | 27   |
| NaN   | 329   | 6   | NaN   | NaN   | 65    | 32     | 25    | 92     | 22   | 11    | 25   | 33   |
| NaN   | 550   | 6   | NaN   | 8     | 38    | 23     | 25    | 56     | 13   | 7     | 25   | 19   |
| NaN   | NaN   | 5   | NaN   | 7     | 31    | 23     | 23    | 47     | 10   | 7     | 23   | 14   |
| NaN   | 1,173 | NaN | NaN   | 71    | 99    | 32     | NaN   | 87     | NaN  | 11    | NaN  | NaN  |
| NaN   | 541   | 6   | 1,595 | NaN   | 60    | 31     | 25    | 89     | 21   | 10    | 25   | 32   |
| NaN   | 273   | NaN | 248   | 8     | 53    | 37     | NaN   | 78     | 17   | 12    | NaN  | 27   |
| NaN   | 230   | 7   | 196   | 7     | 37    | 24     | 24    | 56     | 12   | 7     | 24   | 17   |
| NaN   | 533   | 6   | NaN   | 9     | 59    | 31     | 26    | 90     | 19   | 10    | 26   | 29   |
| NaN   | NaN   | 7   | 205   | 6     | 31    | 32     | 24    | 50     | 10   | 16    | 24   | 17   |
| NaN   | 435   | NaN | 504   | NaN   | 59    | 35     | NaN   | 89     | 23   | NaN   | NaN  | 31   |
| NaN   | 147   | 5   | 137   | 7     | 46    | 29     | 25    | 64     | 15   | 8     | 25   | 20   |
| NaN   | 835   | NaN | NaN   | 44    | 76    | 59     | 41    | 121    | 40   | 36    | 41   | 44   |
| NaN   | 654   | NaN | NaN   | NaN   | 72    | 33     | NaN   | 103    | 24   | 12    | NaN  | 37   |
| NaN   | NaN   | NaN | 8,025 | 1,326 | NaN   | 903    | 652   | 2,878  | NaN  | NaN   | 652  | NaN  |
| NaN   | 334   | 6   | 307   | NaN   | 57    | 37     | 24    | 91     | 21   | 13    | 24   | 33   |
| 2,005 | 2,100 | 9   | 909   | 10    | 34    | 33     | 27    | 52     | 14   | 14    | 27   | 20   |

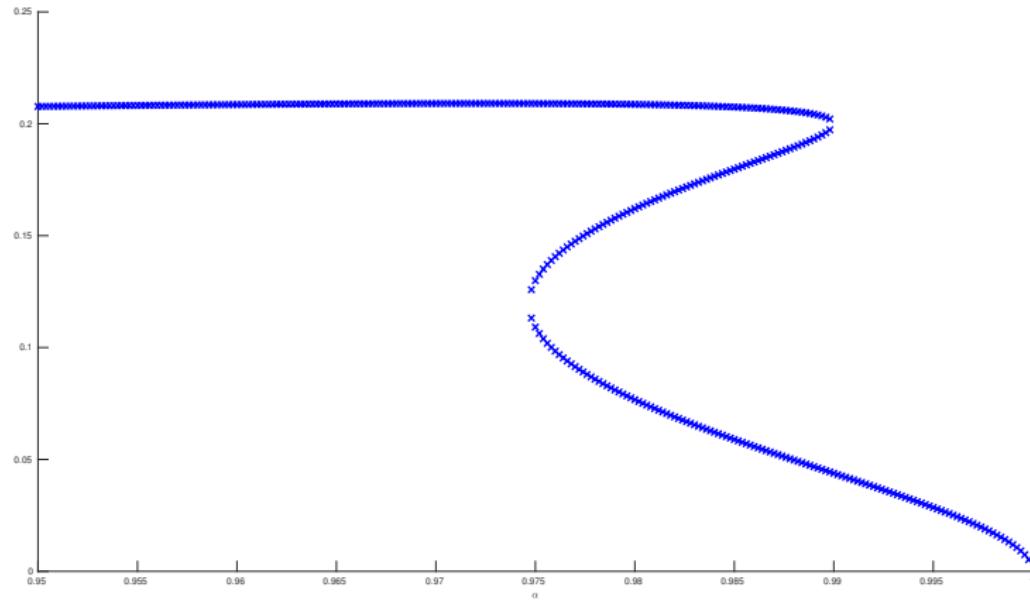
## R6\_3: solution vs. $\alpha$



First entry  $x_1$  of the solution(s) vs. value of parameter  $\alpha$ .

Solutions for smaller values of  $\alpha$  are a poor initial point for  $\alpha = 0.99$ .

## R6\_3: a possible fix



Fix: compute solution for  $\alpha = 0.9999$  first, then use it as starting point to 'go back' to  $\alpha = 0.99$ .

17+8 iterations vs. 1326 iterations.

# Conclusions

- New numerical strategies: Perron-based methods, continuation.
- Can handle all the benchmark problems successfully.
- Insight: consider all solutions (even if we're looking for only one).
- Coming next: tests at larger scale.

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Thanks for your attention!